

Lecture 4: (Advanced?) Decentralized Exchanges

Guillermo Angeris

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Outline

Recap and overview

Examples

Portfolio value function

Conclusion

Previous lecture

- ▶ Discussed decentralized exchanges (DEXs)
- ▶ Focused on a particular implementation: constant function market makers (CFMMs)
- ▶ Talked about some basic notions, such as ‘price’ and ‘portfolio value’

This lecture

- ▶ Recap the previous lecture with concrete examples
- ▶ Discuss some equivalences between portfolio value and trading functions
- ▶ Discuss individual-asset payoffs

Notation recap

- ▶ The notation used here is now (somewhat?) standard
- ▶ $R \in \mathbf{R}_+^n$ are the *reserves*
- ▶ $\varphi : \mathbf{R}_+^n \rightarrow \mathbf{R}$ is the *trading function*
- ▶ $\Delta \in \mathbf{R}_+^n$ is the *tendered basket*
- ▶ $\Lambda \in \mathbf{R}_+^n$ is the *received basket*
- ▶ $0 < \gamma \leq 1$ is the *trading fee*

Interface recap

- ▶ In a CFMM, a `swap(amountIn, amountOut)` accepts if

$$\varphi(R + \gamma\Delta - \Lambda) \geq \varphi(R)$$

where Δ is `amountIn` and Λ is `amountOut`

- ▶ When this happens, reserves are updated as

$$R \leftarrow R + \Delta - \Lambda$$

Arbitrage problem

- ▶ Given external market with price $c \in \mathbf{R}_+^n$, the arbitrage problem is

$$\begin{array}{ll} \text{maximize} & \overbrace{c^T(\Lambda - \Delta)} \\ \text{subject to} & \varphi(R + \gamma\Delta - \Lambda) \geq \varphi(R) \\ & \Delta, \Lambda \geq 0 \end{array} \quad \left. \begin{array}{l} \uparrow \\ \} \end{array} \right\} \leftarrow \begin{array}{l} \text{acceptance} \\ \text{of a CFMM} \end{array}$$

with variables $\Delta, \Lambda \in \mathbf{R}^n$

$\text{OPT} \geq 0$

Arbitrage problem (equiv.)

$$\max. \quad c^T(\underbrace{1 - \Delta}) \quad R' = R + \Delta - \lambda$$

$$\text{s.t.} \quad \varphi(R + \underbrace{\Delta - \lambda}) \geq \varphi(R)$$

$$\Delta, \lambda \geq 0$$

$$\cong \max. \quad c^T(\underbrace{R - R'}) \quad \leftarrow R - R' = \lambda - \Delta$$

$$\text{s.t.} \quad \varphi(R') \geq \varphi(R)$$

$$\cong \max. \quad -c^T R'$$
$$\text{s.t.} \quad \varphi(R') \geq \varphi(R)$$

$$\cong \boxed{\begin{array}{ll} \min. & c^T R' \\ \text{s.t.} & \varphi(R') \geq \varphi(R). \end{array}}$$

Arbitrage problem (equiv.)

- So the final problem is

$$\begin{array}{ll} \text{minimize} & c^T R' \\ \text{subject to} & \varphi(R') \geq \varphi(R) \\ & R' \geq 0 \end{array} \quad \left. \vphantom{\begin{array}{l} \text{minimize} \\ \text{subject to} \end{array}} \right\} V(c)$$

Handwritten red notes:
An arrow points from the text "parameterized by c" to the vector c in the objective function $c^T R'$.
A large red curly brace groups the three lines of the optimization problem, with the label $V(c)$ written to its right.

with variable $R' \in \mathbf{R}^n$

- Call optimal value $V(c)$, the 'portfolio value function'

No-arbitrage conditions

$$\left. \begin{array}{l} \min. c^T R' \\ \text{s.t. } \varphi(R') \geq \varphi(R) \end{array} \right\} \rightarrow$$

$$\mathcal{L}(R') = c^T R' + \lambda(\varphi(R) - \varphi(R')), \quad \lambda \geq 0$$

$$\nabla_{R'} \mathcal{L}(R') = 0 = c - \lambda \nabla \varphi(R') \Rightarrow c = \lambda \nabla \varphi(R')$$

Notes

- ▶ For the most part in this lecture, we will consider $\gamma = 1$ for simplicity
- ▶ This will simplify a lot of results
- ▶ But more general results can be derived

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Linear trading function

$$p \in \mathbb{R}_{++}^n, \quad \varphi(R) = p^T R. \quad 2\text{-asset case}$$

$$\varphi(R + \Delta - \Lambda) \geq \varphi(R)$$

"

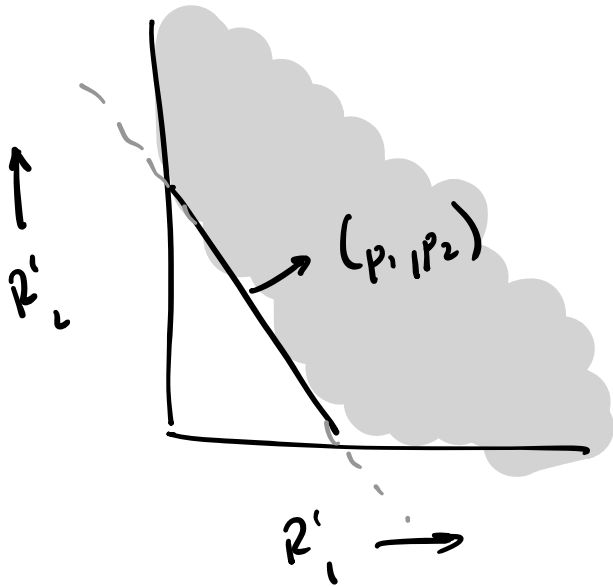
$$\cancel{p^T(R + \Delta - \Lambda)} \geq \cancel{p^T R}$$

$$\boxed{p^T \Delta \geq p^T \Lambda}$$

$$p_1 \Delta_1 + p_2 \Delta_2 \geq p_1 \Lambda_1 + p_2 \Lambda_2$$

$$p_1 \Delta_1 = p_2 \Lambda_2$$

Linear trading function (reachable set, prices)



$$\{R' \mid \varphi(R') \geq \varphi(R)\}.$$

$$\nabla \varphi(R') = (p_1, p_2)$$

$$\varphi(R') = p^\top R'$$

Linear trading function (portfolio value)

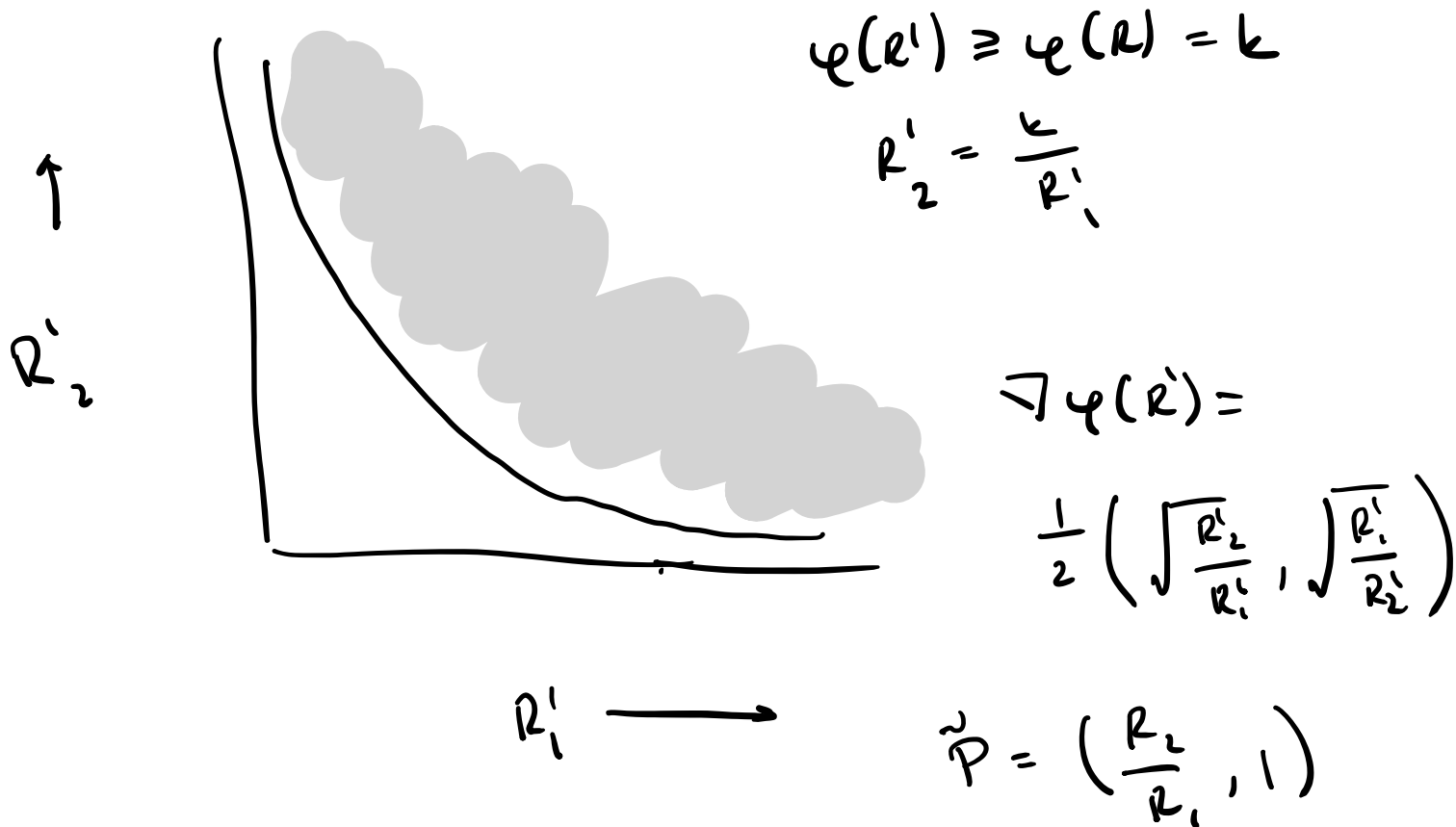
$$\begin{aligned} \min. \quad & c^T R' \\ \text{s.t.} \quad & p^T R' \geq \underbrace{p^T R}_k \Rightarrow g = p_i R'_i \\ & R' \geq 0 \\ & \min. \left(\frac{c}{p} \right)^T g \\ \text{s.t.} \quad & 1^T g \geq k \\ & g \geq 0 \end{aligned}$$

$$V(c) = k \cdot \min_i \left\{ \frac{c_i}{p_i} \right\}.$$

Product trading function

$$\varphi(R_1, R_2) = \sqrt{R_1 R_2}$$

Product trading function (reachable set, prices)



Product trading function (output given input)

Product trading function (portfolio value)

$$c = \lambda \nabla \varphi(k)$$

$$= \frac{\lambda}{2} \left(\underset{\uparrow}{\sqrt{\frac{\bar{R}'_2}{R'_1}}}, \underset{\uparrow}{\sqrt{\frac{\bar{R}'_1}{R'_2}}} \right)$$

$$\frac{c_1}{c_2} = \frac{R'_2}{R'_1} \Rightarrow \boxed{c_1 R'_1 = c_2 R'_2}$$

$$\sqrt{c_1 c_2 k}$$

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Now the question

- ▶ Now, where we were last time
- ▶ Given any portfolio value, can we recover a trading function?

Now the question

- ▶ Not every portfolio value function!
- ▶ By definition must be
 - 1-homogeneous
 - Nondecreasing
 - Concave

$$\begin{array}{ll} \min & c^T R' \\ \text{s.t.} & \varphi(R') \geq \varphi(R) \end{array}$$

Now the question

- ▶ Not every portfolio value function!
- ▶ By definition must be
 - 1-homogeneous
 - Nondecreasing
 - Concave
- ▶ Call these *consistent* portfolio value functions

Now the question

- ▶ Given any **consistent** portfolio value function, can we recover a trading function?

Now the question

- ▶ Given any **consistent** portfolio value function, can we recover a trading function?
- ▶ Yes!

Inverse mapping

$$V(c)$$

$$\tilde{q}(R) = \inf_{c \geq 0} \{ c^T R - V(c) \}.$$

Pf. Show that $\tilde{q}(R)$ has payoff V .

Call payoff of \tilde{q} , \tilde{V} , ① $\tilde{V} \geq V$ ② $\tilde{V} \leq V$

$$\textcircled{1} \underbrace{\left\{ \min_{c \geq 0} c^T R' \right.}_{\text{s.t. } \tilde{q}(R') \geq 0} \left. \right\}$$

if R satisfies then the

$$c^T R' - V(c) \geq \tilde{q}(R') \geq 0$$

$$c^T R' \geq V(c) \Rightarrow \tilde{V} \geq V$$

Inverse mapping (cont.)

$$\textcircled{2} \quad \boxed{R = \nabla V(c)}$$

$$\text{For any } g: \quad V(g) \leq \overbrace{\nabla V(c)^T}^R (g - c) + V(c)$$

$$\textcircled{a} \quad c^T R = V(c)$$

$$\textcircled{b} \quad R \text{ is feasible for } \tilde{V}$$

$$\Rightarrow \tilde{V} \leq V \Rightarrow V = \tilde{V}$$

What does this give us?

- ▶ We can construct complicated, interesting payoffs
- ▶ One example: we can construct fully decentralized options !
- ▶ These options do not require counterparties
- ▶ Only rely on people arbitraging the protocol

Second view on CFMMs

- ▶ We can view CFMMs as fully-decentralized portfolio management
- ▶ Only requires someone being able to arbitrage portfolio
- ▶ Risk-return tradeoff is specified by V
- ▶ Can construct a φ which gives this tradeoff

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Final points

- ▶ Number of other interesting results in CFMMs
- ▶ Including ways of constructing very general payoffs
- ▶ Could be a special topic...

Next lecture

- ▶ Things that can only be done on the blockchain
- ▶ Flashloans, atomic arbitrage, sandwiching...
- ▶ And more **spicy** topics!